

Seminář 5

The system $\lambda\mu$

- It is the system of *recursive types*.
- The recursive types come together with an equivalence relation \approx .
- The type assignment rules consist of the rules of $\lambda \rightarrow$ and the following rule

$$\frac{\Gamma \vdash M : \sigma \quad \sigma \approx \sigma'}{\Gamma \vdash M : \sigma'}$$

The set $\mathbf{T} = \text{Type}(\lambda\mu)$, trees of types of $\lambda\mu$.

(i) The set of types of $\lambda\mu$, $\mathbf{T} = \text{Type}(\lambda\mu)$, is defined by the following abstract grammar.

$$\mathbf{T} = \mathbf{V} \mid \mathbf{T} \rightarrow \mathbf{T} \mid \mu \mathbf{V}. \mathbf{T}$$

where \mathbf{V} is the set of type variables.

(ii) Let $\sigma \in \mathbf{T}$ be a type. The *tree of* σ , $T(\sigma)$ is defined by induction on the structure of σ as follows:

$$T(\alpha) = \alpha \quad \text{if } \alpha \text{ is a type variable}$$

$$T(\sigma \rightarrow \tau) = \begin{array}{c} \rightarrow \\ \swarrow \quad \searrow \\ T(\sigma) \quad T(\tau) \end{array}$$

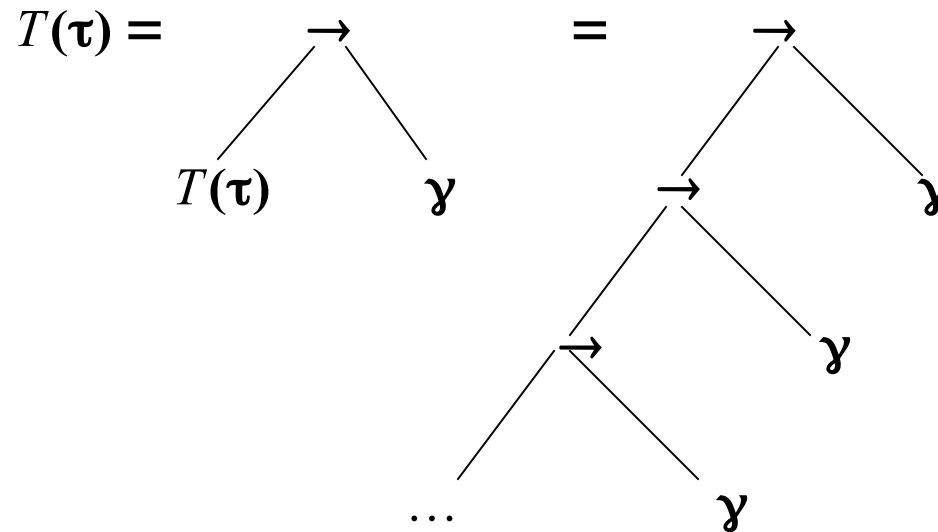
$$T(\mu\alpha.\sigma) = \begin{cases} \perp & \text{If } \sigma \equiv \mu\beta_1 \dots \beta_n.\alpha \text{ for} \\ & \text{some } n \geq 0 \\ T(\sigma[\alpha := \mu\alpha.\sigma]) & \text{else} \end{cases}$$

(iii) The equivalence relation \approx on trees is defined as follows:

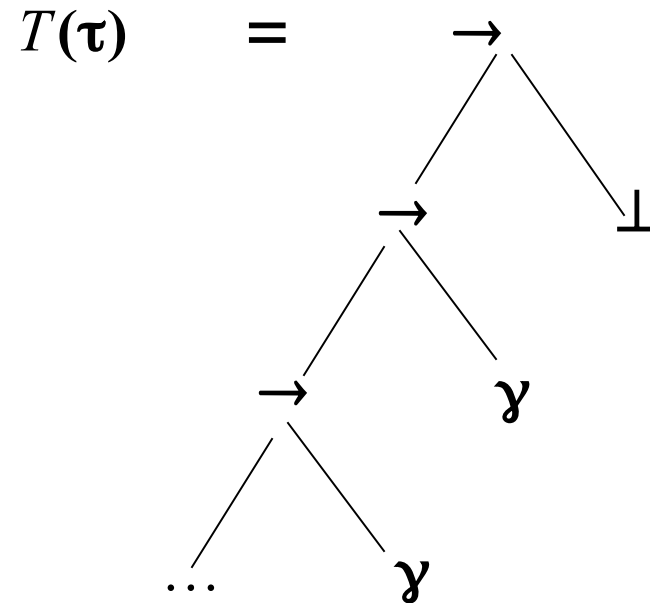
$$\sigma \approx \tau \Leftrightarrow T(\sigma) = T(\tau)$$

Exercises.

(a) Assume $\tau \equiv \mu\alpha.\alpha \rightarrow \gamma$, then



(b) Assume $\tau \equiv (\mu\alpha.\alpha \rightarrow \gamma) \rightarrow \mu\delta\mu\beta.\beta$, then



(c) $(\mu\alpha.\alpha \rightarrow \gamma) \approx (\mu\alpha(\alpha \rightarrow \gamma) \rightarrow \gamma).$

(d) $\mu\alpha.\sigma \approx \sigma[\alpha := \mu\alpha.\sigma]$ for all σ , even if $\sigma \equiv \mu\bar{\beta}.\alpha.$

$\lambda\mu$

(start rule)	$\frac{(x:\sigma)\in\Gamma}{\Gamma -x:\sigma}$
(\rightarrow -elimination - \leftarrow)	$\frac{\Gamma -M:(\sigma\rightarrow\tau) \quad \Gamma -N:\sigma}{\Gamma -(MN):\tau}$
(\rightarrow -introduction - \leftarrow)	$\frac{\Gamma, x:\sigma -M:\tau}{\Gamma -(\lambda x.M):(\sigma\rightarrow\tau)}$
(\approx -rule)	$\frac{\Gamma -M:\sigma \quad \sigma\approx\tau}{\Gamma -M:\tau}$

$\neg \vdash K : \alpha \rightarrow \alpha$

$\neg \vdash (\lambda x. (\lambda y. x)) : \alpha \rightarrow \alpha$